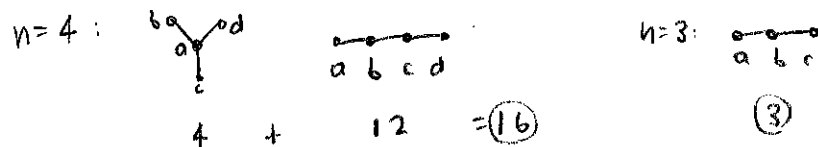


Lecture 15
10.17.13

Trees

A tree on $[n]$ is a graph on vertices $1 \dots n$ which is connected and has no cycles. $t(n) = \#$ of them



A rooted tree is a tree with a chosen vertex called the root. $r(n) = \#$ of them RT = comb. class.

$$r(n) = n t(n).$$

Note: A rooted tree is



a root and a "rooted forest"

• RT = Atom \star RF

• RF = Set (RT)

So

$$R(z) = ze^{R(z)}$$

One interpretation: $R(z) = z e^{z e^{z e^{z e^{\dots}}}}$

Another: $z = R(z) e^{-R(z)}$

↓

$R(z) = (z e^{-z})^{\leftarrow 1}$ (compositional inverse)

How do I find the coeffs. of a comp inverse?

Fact

A power series $f(x) = a_1 x + a_2 x^2 + \dots \in \mathbb{R}[[x]]$ has a compositional inverse $f^{\leftarrow}(x)$ iff $a_1 \neq 0$, in which case $f^{\leftarrow}(f(x)) = f(f^{\leftarrow}(x)) = x$

Lagrange Inversion Formula: If $f^{\leftarrow}(x)$ exists,

$$n [x^n] f^{\leftarrow}(x)^k = k [x^{n-k}] \left(\frac{x}{f(x)} \right)^n$$

Then

$$\begin{aligned} \frac{n r(n)}{n!} &= n [x^n] R(x) \\ &= n [x^n] (x e^{-x})^{\leftarrow 1} \\ &= [x^{n-1}] \left(\frac{x}{x e^{-x}} \right)^n \\ &= [x^{n-1}] e^{nx} \\ &= \frac{n^{n-1}}{(n-1)!} \end{aligned}$$

Sylvester-Carley

$$\Rightarrow \begin{cases} r(n) = n^{n-1} \\ t(n) = n^{n-2} \end{cases}$$

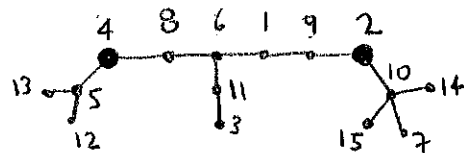
(63)

How to prove Lagrange inversion?

- analysis
- combinatorially, using comb. of trees!

Such a nice formula deserves a bijective proof!

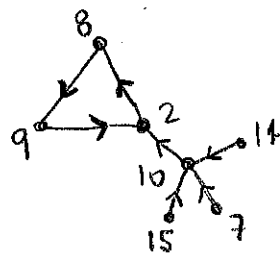
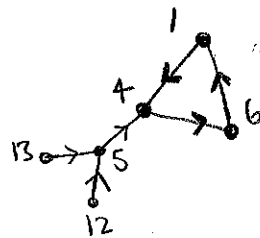
- ① Look up "Prüfer code"
- ② Need to show there are n^n "trees with skeleton"



$$\begin{pmatrix} 1 & 2 & 4 & 6 & 8 & 9 \\ 4 & 8 & 6 & 1 & 9 & 2 \end{pmatrix}$$

=

$$(146)(289)$$



Turn the permutation of the skeleton into cycle notation, draw the graph of the permutation, "rehang" the trees hanging from the skeleton, and point all arrows towards the cycles.

The result is the graph of a function $f: [n] \rightarrow [n]$ and this mapping is a bijection.

(64)

Parking function

There are n cars $C_1 \dots C_n$ trying to park on the n spots $1 \dots n$ of a one way street.

Car C_i has a preferred spot a_i ; it takes that one, or the first available after it.

$n \dots 3 \ 1$

$1 \ 2 \ 3 \dots n$

Ex: $a = (3, 2, 3, 1, 2) \Rightarrow$ $4 \ 2 \ 1 \ 3 \ 5$

$a = (4, 3, 3, 4, 1) \Rightarrow$ $1 \ 2 \ 3 \ 4 \dots ?$

a is a parking function if all cars can park.

Prop a is a parking function



it contains $\geq i$ numbers $\leq i$ for all i .

Eqv: Rearranging it to get $b_1 \leq b_2 \leq \dots \leq b_n$, I need $b_i \leq i$ for all i .

\Downarrow If not, there are $> n-i$ numbers $> i$
 $\Rightarrow > n-i$ cars fighting for $n-i$ spots

\Uparrow Exercise

Ex: $n=3$:

111					
112	121	211			
113	131	311			
122	212	221			
123	132	213	231	312	321

16 parking fns

5 classes

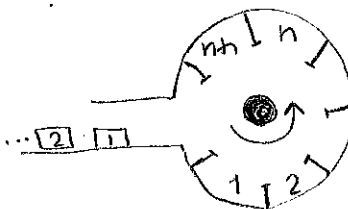
↖
 Catalan #.
 Why?

(65)

Theorem There are $(n+1)^n$ parking functions of length n

Proof

Put them in a "rampoi" with $n+1$ spots instead:



- If a car can't find a spot, it keeps circling around
- Now $n+1$ is a possible preferred spot

$a = (4, 3, 3, 4, 1) \rightarrow$



Now all cars can park and exactly one spot is left empty. Note:

- if (a_1, a_2, \dots, a_n) leaves i empty
 $(a_1+1, a_2+1, \dots, a_n+1)$ leaves $i+1$ empty (mod $n+1$)
- then there are $\frac{1}{n+1} \cdot (n+1)^n$ preference fns that leave $n+1$ empty
- these are precisely the parking functions!

Q Bijection?

(parking fns. of size n) \leftrightarrow (trees on $[n+1]$)

• Think
 • Google (66)