

Counting walks in graphs (to count reduced words)

G-digraph

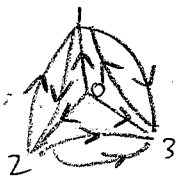
$a_{uv} = \#$ of edges $u \rightarrow v$

↑
adjacency matrix

$P_{uv}(n) = \#$ of paths

$u \rightarrow \dots \rightarrow v$
n edges

Ex
automaton for $\infty \Delta$



$$A = \begin{matrix} & \begin{matrix} 0 & 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix} \end{matrix}$$

Goal: Find $P_{uv}(n)$.

Transfer-Matrix Method (statistical mechanics / Markov chains in prob.)

Note:

$$P_{uv}(n) = \sum_w a_{uw} P_{wv}(n-1)$$

$$P(n) = A \quad P(n-1) = A \quad \dots \quad P(0) = I$$

Theorem Fix u, v .

$$\sum_{n \geq 0} P_{uv}(n) x^n = (-1)^{u+v} \frac{\det(I - XA; j, i)}{\det(I - XA)}$$

← del. row j, col i

Pf. $\sum_{n \geq 0} P_{uv}(n) x^n = \sum_{n \geq 0} (A^n)_{uv} x^n$
 $= \left(\sum_{n \geq 0} A^n x^n \right)_{uv} = (I - Ax)^{-1}_{uv}$ as claimed. \square

In Ex, $\sum_{n \geq 0} P_{02}(n) x^n = \det \begin{vmatrix} -x & -x & -x \\ 0 & 1-x & -x \\ 0 & -x & 1-x \end{vmatrix} = \frac{x^2 - x^3 + x^2 + x^3 + x^2 + x}{1 + x^3 + x^3 - x^2 - x^2 - x^2}$
 $\det \begin{vmatrix} 1 & -x & -x & -x \\ 0 & 1-x & -x & -x \\ 0 & -x & 1-x & -x \\ 0 & -x & -x & 1 \end{vmatrix} = \frac{x + 2x^2 + x^3}{1 - 3x^2 - 2x^3}$
 $= \frac{x(1+x)^2}{(1+x)^2(1-2x)} = \frac{x}{1-2x}$

So if $r_n = \#$ of reduced words of length n ,

$$r_n = \begin{cases} 1 & n=0 \\ 3P_{02}(n) & n \geq 1 \end{cases}$$

$$\sum_{n \geq 0} r_n x^n = 1 + \frac{3x}{1-2x} = 1 + \sum_{n \geq 0} 3 \cdot 2^n x^n$$

$$r_n = 3 \cdot 2^n \quad n \geq 1$$

Not too impressive, but this is $\infty \Delta$!

Corollary $R_{(w,s)}(x) = \sum_{n \geq 0} r_n x^n$ is a rational function for all (w,s)

Often this is better than a formula for r_n .

$$R_{\Delta}(x) = \frac{(1+2x)(1+x^2)}{(1-x)(1-2x^2)} ; r_n = \underbrace{-(\sqrt{2})^n}_{\uparrow} + \underbrace{-(-\sqrt{2})^n}_{\uparrow} + \dots$$