

Now need  $l(ws_i) > l(w) \Rightarrow w\alpha_i > 0$

Induct on  $l(w)$ . ( $l(w)=0$  is obvious)

Take  $s$  with  $ws < w$ .

Let  $J = \{s, s_i\}$  and use  $W = W^J W_J$

$\Rightarrow W = W^J W_J$

- $l(w) = l(w^J) + l(w_J)$
- $w^J = \min$  in coset  $wW_J$

•  $W\alpha_i = W^J W_J \alpha_i$

$\rightarrow$  Note  $l(W_J s_i) > l(W_J)$  since

$$\begin{array}{ccc} Ws_i & = & W^J W_J s_i \\ \downarrow & & \downarrow \\ W & & W^J W_J \end{array} \quad \begin{array}{l} l(ws_i) = l(w^J) + l(w_J s_i) \\ \downarrow \\ l(w) = l(w^J) + l(w_J) \end{array}$$

so by dihedral case

$$W_J \alpha_i = a\alpha_i + b\alpha \quad a, b \geq 0$$

$\rightarrow$  So  $w\alpha_i = w^J(a\alpha_i + b\alpha)$

Is  $w^J \alpha_i > 0$ ? Is  $w^J \alpha > 0$ ?

By induction, enough to show

$$l(W^J s_i) > l(W^J) \quad l(W^J s) > l(W^J)$$

But these follow since

$$W^J s_i = W W_J^{-1} s_i \in w W_J$$

$$W^J s \in w W_J \quad \square$$

Note.

$$\begin{array}{l} l(ws) > l(w^J) \\ \text{Since } w \notin W^J \\ l(ws) < l(w) \end{array}$$

Corollary The geometric rep'n is faithful

(injective)

Pf.  $\sigma: W \rightarrow GL(Y)$

Sup  $\sigma(w) = \text{id}$ .

$\Rightarrow W\alpha_i = \alpha_i > 0$  for all  $\alpha_i$

$\Rightarrow l(ws_i) > l(w)$  for all  $s_i$

$\Rightarrow w = e. \quad \square$

Prop.

(a) Let  $s_i \in S$ . Then  $s_i$  sends  $\alpha_i$  to  $-\alpha_i$  and permutes the other positive roots

(b) Let  $w \in W$ . Then

$l(w) = \#$  of pos. roots that  $w$  sends to neg. roots

Pf

(a) Clearly  $s_i(\alpha_i) = -\alpha_i$ . Now let  $\alpha \in \Phi^+$   
 $\alpha \neq -\alpha_i$

Write

$$\alpha = \sum_j c_j \alpha_j \quad \text{some } c_j > 0 \quad j \neq i$$

$$s_i(\alpha) = \alpha - 2 \langle \alpha, \alpha_i \rangle \alpha_i$$

$\uparrow$  coeff of  $\alpha_j$  is still  $c_j > 0$

$\Rightarrow s_i(\alpha) > 0$ .

So  $s_i: \Phi^+ \setminus \{\alpha_i\} \rightarrow \Phi^+ \setminus \{\alpha_i\}$

$$s_i s_j: \alpha = s_i(s_j \alpha) \quad \text{Inj: } s_i \alpha = s_i \beta \Rightarrow s_i s_i \alpha = s_i s_i \beta$$

(b) Induct on  $l(w)$ .  $l(w)=1$  is ok by (a).

Sup true for length  $< l(w)$ .

Take  $s_i$  with  $l(ws_i) < l(w)$

$$(w\alpha_i < 0)$$

$\Rightarrow l(ws_i)$  roots  $\alpha > 0$  with  $ws_i\alpha < 0$

$\rightarrow l(ws_i)$  roots  $s_i\alpha > 0$  with  $w(s_i\alpha) < 0$   
 $\hookrightarrow$  also not  $\alpha_i$

$\rightarrow$  another one:  $w\alpha_i < 0$

$\rightarrow$  Clearly no others.  $\Rightarrow l(ws_i) \neq l(w)$

Roots, alg. reflections, geom. reflection

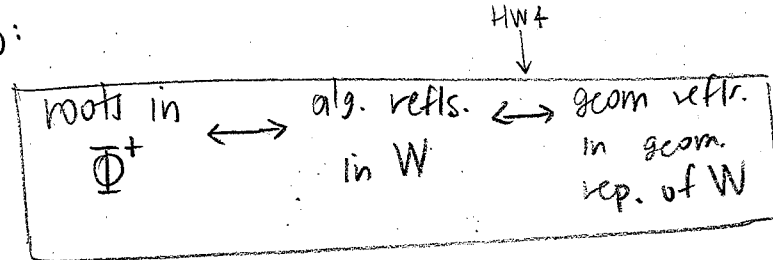
Consider how an alg. reflection  $t = ws_i w^{-1}$  acts:

$$\begin{aligned} tv &= ws_i w^{-1}(v) \\ &= w[w^{-1}(v) - 2\langle w^{-1}(v), \alpha_i \rangle \alpha_i] \\ &= v - 2\langle w^{-1}(v), \alpha_i \rangle w\alpha_i \\ &= v - 2\langle v, w\alpha_i \rangle w\alpha_i \end{aligned}$$

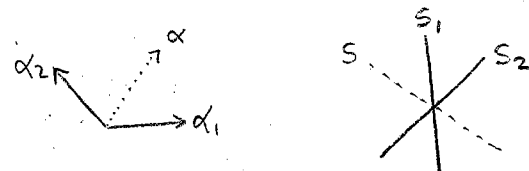
action of alg refl:  $ws_i w^{-1}$  = geom refl. by root  $w\alpha_i$   
 $\parallel$  reflection  $\parallel$  root

Lemma.  
 $\langle wv_1, wv_2 \rangle = \langle v_1, v_2 \rangle$   
PF.  
 Compute for each  $i$ ,  
 $\langle s_i v_1, s_i v_2 \rangle = \langle v_1, v_2 \rangle$

So:



Ex:  $S_3$



Prop. Let  $w \in W$ ,  $t \in T$  reflection,  $\alpha = \text{corr. root}$

$$l(wt) > l(w) \Rightarrow w\alpha > 0$$

$$l(wt) < l(w) \Rightarrow w\alpha < 0$$

PF As in the case  $t \in S$ , enough to show one

Sup  $l(wt) < l(w)$

$$w = s_1 s_2 \dots s_k$$

$$t = \underbrace{s_k \dots s_i}_{U} s_i \underbrace{s_i \dots s_k}_{U^{-1}}$$

$$\downarrow$$

$$\alpha = U\alpha_i$$

$$w\alpha = wU\alpha_i = s_1 s_2 \dots s_i \alpha_i < 0$$

since  $s_i$  shortens  $s_1 \dots s_i$