

The proposition gives us a "projection" map

$$P^J : W \rightarrow W^J$$

$$w \mapsto w^J \text{ such that } w = w^J w_J \quad (w^J \in W^J, w_J \in W_J)$$

Prop  $P^J$  is order-preserving

$$\text{PF Need: } w_1 \leq w_2 \Rightarrow w_1^J \leq w_2^J$$

Induct on  $\ell(w_2)$ .  $\ell(w_2) = \ell(w_1)$  is ok.

Suppose the for length  $< \ell(w_2)$ .

$$w_1^J \leq w_1 \leq w_2 \text{ so done unless } w_2^J \neq w_2$$

$$\Rightarrow \text{Assume } w_2^J \neq w_2 \Rightarrow w_2 = w_2^J (w_2)_J$$

$$\ell(w_2) = \ell(w_2^J) + \ell((w_2)_J)$$

let  $j = \underset{\text{(some)}}{\underset{1}{\overset{n}{\exists}}} \text{ letter of } (w_2)_J \Rightarrow w_2 j \leq w_2$

$$\begin{array}{ccc} w_2 & & \\ \swarrow w_2 j & \searrow w_1^J j & \\ w_1^J & & \end{array}$$

$$\rightarrow w_1^J \leq w_2 j$$

↓ induction

$$(w_1^J)^J \leq (w_2 j)^J$$

$$w_1^J \leq w_2^J$$

$$? w_2 j = \underbrace{w_2^J}_{\text{B}} (w_2)_J j$$

An enumerative application: length-generating fns

The Poincaré series of  $W$  is

$$W(q) = \sum_{w \in W} q^{\ell(w)}$$

$$\text{Ex. } S_3(q) = 1 + 2q + 2q^2 + q^3$$

$$S_4(q) = 1 + 3q + 5q^2 + 6q^3 + 5q^4 + 3q^5 + q^6$$

Prop  $W(q)$  is symmetric for  $W$  finite

PF  $w \mapsto w w_0$  is an antiautomorphism.

$$W\left(\frac{1}{q}\right) = \frac{W(q)}{q^{1/2}}$$

Lemma  $W(q) = W^J(q) W_J(q)$

$$\text{PF RHS} = \sum_{w^J \in W^J} q^{\ell(w^J)} \sum_{w_J \in W_J} q^{\ell(w_J)}$$

$$= \sum_{\substack{w^J \in W^J \\ w_J \in W_J}} q^{\ell(w^J w_J)} = \text{LHS}$$

So if I can compute  $W^J(q)$  I can recursively compute  $W(q)$

Q What about  $w \mapsto w_J$ ?

Ex  $W = S_n$

$$J = \{s_1, \dots, s_{n-2}\}$$

$S_n^J$  - elt whose reduced expression

cannot end in  $s_1, \dots, s_{n-2}$

$$= \{e, s_{n-1}, s_{n-2}s_{n-1}, \dots, s_1 \dots s_m\}$$

$$S_n^J(q) = 1 + q + q^2 + \dots + q^{m-1}$$

$$\rightarrow S_n(q) = S_n^J(q)(S_n)_J(q)$$

$$= (1 + q + \dots + q^{m-1}) S_m(q)$$

$$\rightarrow \boxed{\text{Thm. } S_n(q) = (1+q)(1+q+q^2) \dots (1+q+\dots+q^{m-1})}$$

$$= [n]! = [n][n-1] \dots [2][1]$$

$$\text{where } [i]_q = 1 + q + \dots + q^{i-1}$$

In fact,

$$W = W_1 \times \dots \times W_n \Rightarrow W(q) = W_1(q) \cdots W_n(q) \quad (\text{easy})$$

and

Thm If  $(W, S)$  is finite, irreducible, and  $n = |S|$ ,  
there are "exponents"  $e_1, \dots, e_n \in \mathbb{N}$  such that

$$W(q) = \prod_{i=1}^n [e_i + 1]_q$$

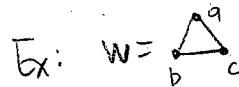
$$\text{so } |W| = \prod_{i=1}^n (e_i + 1) \text{ and } |\Pi| = \ell(W_0) = \sum_{i=1}^n e_i$$

Pf: Conceptually  $\rightarrow$  invariant theory

Manually  $\rightarrow$  inductively using the classification  
of finite Coxeter groups.

$\text{Thm}$ $\sum_{J \subseteq S} \frac{(-1)^{ J }}{W_J(q)} = \begin{cases} \frac{q^{\ell(W_0)}}{W(q)} & W \text{ finite} \\ 0 & W \text{ infinite} \end{cases}$
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Note: This allows recursive computation of  $W(q)$ .



$$\begin{array}{c|c|c}
J & W_J(q) \\
\hline
\emptyset & 1 \\
a, b, c & 1+q \\
ab, ac, bc & (1+q+q^2)(1+q) \\
abc & ?
\end{array}$$

$$1 - \frac{3}{1+q} + \frac{3}{(1+q)(1+q+q^2)} - \frac{1}{W(q)} = 0 \quad ab, ac, bc \quad (1+q+q^2)(1+q)$$

$$W(q) = \frac{1+2q+2q^2+q^3}{1-q-q^2+q^3}$$

for infinite  $W$ ,  
 $W(1/q) = \pm W(q)$   
also "symmetric"

PF Mult by  $W(q)$ .

$$\sum_{J \subseteq S} (-1)^{|J|} W^J(q) = \sum_{J \subseteq S} \sum_{w \in W^J} (-1)^{|J|} q^{\ell(w)}$$

$$= \sum_{w \in W} q^{\ell(w)} \sum_{J \subseteq S} (-1)^{|J|} = \sum_{w \in W} q^{\ell(w)} \sum_{J \subseteq S \setminus D(w)} (-1)^{|J|}$$

$$= \sum_{w \in W} q^{\ell(w)} (1 + (-1))^{S \setminus D(w)}$$

$$= \sum_{\substack{w \in W \\ D(w) = S}} q^{\ell(w)} = \begin{cases} q^{\ell(W_0)} & W \text{ finite} \\ 0 & W \text{ infinite} \end{cases}$$

$\uparrow$  descrb. of  $w$   
= possible last  
letters of reduced  
words for  $w$   
= s's that  
shorten  $w$