

The "long word"

Lemma For every $u, v \in W$ there is $w \in W$
with $u \leq w, v \leq w$

Pf Induct on $l(u) + l(v)$

• case 0: \checkmark

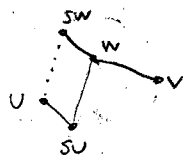
• case $n-1 \rightarrow$ case n :

Take s with $l(su) = l(u) + 1$

Take $su \leq w, v \leq w$

If $w < sw$:

lift to $u \leq sw$
 $v \leq sw$



If $w > sw$:

lift to $u \leq w$
 $v \leq w$



Corollary A finite Coxeter group has a
unique max elt w_0 , the "long word".

Converse: If there is $w_0 \in W$ with $sw_0 < w_0$
for all $s \in S$, then W is finite, w_0 is its max

Pf: $u \leq w_0$ follows by induction on $l(u)$.

W is finite since all of W can be realized
as subwords of a reduced word for w_0 \square

For $W = S_n$, $w_0 = n(n-1)\dots 321$.

Properties

(i) $w_0^2 = e$

(ii) $l(w w_0) = l(w_0) - l(w)$

(iii) $T_L(w w_0) = T \setminus T_L(w)$

(iv) $l(w_0) = |T|$

Recall:
 $T_L(u) = \{ \text{reflections } t \mid$
 $t u < u \}$
 $l(u) = |T_L(u)|$

Pf

(i) Since $l(w_0^{-1}) = l(w_0)$, $w_0^{-1} \leq w_0 \Rightarrow w_0^{-1} = w_0$

(ii) \geq : by triangle inequality

\leq : Induct down on $l(w)$.

$l(w) = l(w_0)$: $w = w_0$ \checkmark

$l(w) = n - m$:

Take $sw > w$

$l(sw w_0) \leq l(w_0) - l(sw)$

$l(w w_0) - 1 \leq l(w_0) - l(w) - 1$ \checkmark

(iii) $t \in T_L(w w_0) \Leftrightarrow t w w_0 < w w_0$

$\Leftrightarrow l(t w w_0) < l(w w_0)$

$\Leftrightarrow l(t w) > l(w)$

$\Leftrightarrow t w > w$

$\Leftrightarrow t \notin T_L(w)$

(iv) Apply (iii) to $w = e$ \square